

Secure Multi-Execution

Dominique Devriese Frank Piessens

K.U.Leuven

May 14, 2010

Outline

Secure Multi-Execution

Introduction

Informal Overview

Formal Properties

Experimental Results

Conclusion

Introduction

- ▶ Information Flow Analysis has received much attention:
 - ▶ Static analysis methods:
From Denning to JFlow/JIF and FlowCaml
But:
 - ▶ Substantial Programmer Effort
 - ▶ In general undecidable statically
 - ▶ Hard to handle exceptions, parallelism, timing covert channel
 - ▶ Dynamic methods:
Many practical but unsound methods, some sound and somewhat practical methods
But:
 - ▶ Some use cases require sound methods (e.g. web page scripts)
 - ▶ No existing monitor can precisely enforce non-interference
 - ▶ Hard to handle exceptions, parallelism, timing covert channel

Secure Multi-Execution

Secure Multi-Execution

- ▶ A novel dynamic enforcement technique for non-interference
- ▶ Nice theoretical properties
 - ▶ Strong soundness guarantee
 - ▶ The first (afawk) sound and precise enforcement method
- ▶ Practical in some scenario's
 - ▶ Performance measurements
 - ▶ Browser implementation possible?

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Information Flow Analysis

```
1 var text = document.getElementById
2     ('email-input').text;
3 var abc = 0;
4 if(text.indexOf('abc')!=-1) { abc = 1 };
5 var url = 'http://example.com/img.jpg'
6     + '?t=' + escape(text) + abc;
7 document.getElementById('banner-img')
8     .src = url;
```

Information Flow Analysis

```
1 var text = document.getElementById
2     ('email-input').text; ← Input at level H
3 var abc = 0;
4 if(text.indexOf('abc')!=-1) { abc = 1 };
5 var url = 'http://example.com/img.jpg'
6     + '?t=' + escape(text) + abc;
7 document.getElementById('banner-img')
8     .src = url; ← Output at level L
```

Timing Covert Channel

```
1 function time(f) {
2   var t = new Date().getTime();
3   f();
4   return new Date().getTime() - t;
5 }
6 function f() {
7   if(abc != 0) {
8     for(var i = 0; i < 10000; ++i) {}
9   }
10 }
11 var abcLo = 0
12 if(time(f) > 10) {
13   abcLo = 1;
14 }
```


Termination Covert Channel

```
1 while(abc == 0) {}  
2 img.url = 'http://example.com/img.jpg';
```

Secure Multi-Execution

L

```

1 var t = (...).text
2 var abc = 0
3 if(t.indexOf('abc')!=-1)
4   { abc = 1 }
5 var url = baseUrl + '?t='
6   + escape(t) + abc
7 (...).src = url

```

H

```

1 var t = (...).text
2 var abc = 0
3 if(t.indexOf('abc')!=-1)
4   { abc = 1 }
5 var url = baseUrl + '?t='
6   + escape(t) + abc
7 (...).src = url

```

Secure Multi-Execution

L

```

1 var t = (...).text undefined
2 var abc = 0
3 if(t.indexOf('abc')!=-1)
4   { abc = 1 }
5 var url = baseUrl + '?t='
6   + escape(t) + abc
7 (...).src = url

```

H

```

1 var t = (...).text
2 var abc = 0
3 if(t.indexOf('abc')!=-1)
4   { abc = 1 }
5 var url = baseUrl + '?t='
6   + escape(t) + abc
7 (...).src = url

```

Secure Multi-Execution

L

```

1 var t = (...).text undefined
2 var abc = 0
3 if(t.indexOf('abc')!=-1)
4   { abc = 1 }
5 var url = baseUrl + '?t='
6   + escape(t) + abc
7 (...).src = url

```

H

```

1 var t = (...).text
2 var abc = 0
3 if(t.indexOf('abc')!=-1)
4   { abc = 1 }
5 var url = baseUrl + '?t='
6   + escape(t) + abc
7 (...).src = url

```

Input Side Effects

L

```

1 var t = (...).text undefined
2 var c = window.confirm
3 if( c("Send e-mail?" )
4   { (...) }
5 var abc = 0
6 if(t.indexOf('abc')!=-1)
7   { abc = 1 }
8 var url = baseUrl + '?t='
9   + escape(t) + abc
10 (...).src = url

```

H

```

1 var t = (...).text
2 var c = window.confirm
3 if( c("Send e-mail?" )
4   { (...) }
5 var abc = 0
6 if(t.indexOf('abc')!=-1)
7   { abc = 1 }
8 var url = baseUrl + '?t='
9   + escape(t) + abc
10 (...).src = url

```

Input Side Effects

L

```

1 var t = (...).text
2 var c = window.confirm
3 if( c("Send e-mail?") )
4   { (...) }
5 var abc = 0
6 if(t.indexOf('abc')!=-1)
7   { abc = 1 }
8 var url = baseUrl + '?t='
9   + escape(t) + abc
10 (...).src = url

```

undefined

H

```

1 var t = (...).text
2 var c = window.confirm
3 if(c("Send e-mail?") )
4   { (...) }
5 var abc = 0
6 if(t.indexOf('abc')!=-1)
7   { abc = 1 }
8 var url = baseUrl + '?t='
9   + escape(t) + abc
10 (...).src = url

```

Secure Multi-Execution

Properties

- ▶ “Obviously” sound:
 - Only execution at high level can see the real high inputs
 - Only execution at low level can produce low outputs
- ▶ “Obviously” precise:
 - If a program is non-interferent, then changing high inputs in low executions will not change their low behaviour

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Non-interference

A formalization of information flow policies

Assume:

- ▶ sets of input channels \mathcal{C}_i , output channels \mathcal{C}_o
- ▶ security level lattice \mathcal{L}
- ▶ $\sigma_{in} : \mathcal{C}_i \rightarrow \mathcal{L}$, $\sigma_{out} : \mathcal{C}_o \rightarrow \mathcal{L}$
- ▶ inputs $I : \mathcal{C}_i \rightarrow (\mathbb{N} \rightarrow Int)$, outputs $O : \mathcal{C}_o \rightarrow List[Int]$
- ▶ $I =_l I'$ iff $I(c_i) = I'(c_i)$ for all c_i such that $\sigma_{in}(c_i) \leq l$

Definition

A program P is (termination-insensitively) *non-interferent* if for all security levels l and inputs $I =_l I'$, where P terminates for I and I' with outputs O and O' , we have that $O =_l O'$.

Soundness

Definition (Strong non-interference)

A program P is *timing-sensitively non-interferent* or *strongly non-interferent* with relation to a given semantics \hookrightarrow^* if for all security levels $l \in \mathcal{L}$, for all $n \geq 0$, for all program inputs I and I' such that $I =_l I'$ holds that if

$$(P, I) \hookrightarrow^n (p, O) ,$$

then

$$(P, I') \hookrightarrow^n (p', O') ,$$

and $p' =_l p$ and $O' =_l O$.

Theorem (Soundness of Secure Multi-Execution)

Any program P is strongly non-interferent under secure multi-execution, using the $\text{select}_{\text{lowprio}}$ scheduler function.

Precision

Theorem (Precision of Secure Multi-Execution)

*Suppose we have a (termination-sensitively) non-interferent program P .
Suppose that*

$$(P, I) \rightarrow^* (p, O)$$

(terminates) for some I , p and O . Then

$$(P, I) \Rightarrow^* (p, O) .$$

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Experimental Results

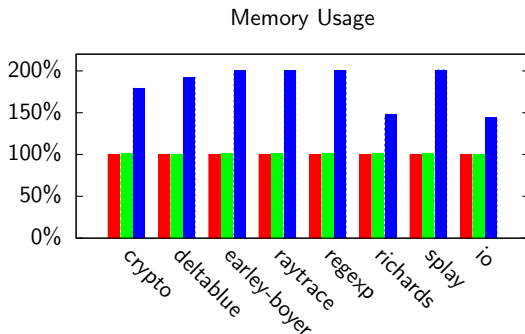
- ▶ Spidermonkey Javascript engine, no real browser
- ▶ 2 security levels
- ▶ Dual-core PC
- ▶ 3 types of execution:
 - ▶ Standard execution
 - ▶ Serial Multi-Execution
 - ▶ Parallel Multi-Execution
- ▶ Benchmarks:
 - ▶ Google Chrome V8 Benchmark Suite: crypto, deltablue, earley-boyer, raytrace, regexp, richards, splay
 - ▶ io: model I/O functions: hi_input, hi_output, lo_input, lo_output: some calculations + I/O at different security levels

IO Benchmark




```
1 for (var i = 0; i < 100; ++i) {
2   var test = 0;
3   for (var j = 0; j < 10000; ++j) {
4     test += j;
5   }
6   if (i % 10 == 0) {
7     var hi_in = hi_input();
8     var lo_in = lo_input();
9     lo_output("#" + i + ". lo_in: '"
10      + lo_in + "'. hi_in is: '"
11      + hi_in + "'");
12     hi_output("#" + i + ". hi_in: '"
13      + hi_in + "'. lo_in is: '"
14      + lo_in + "'");
15   }
16 }
```

Experimental Results: Memory Usage

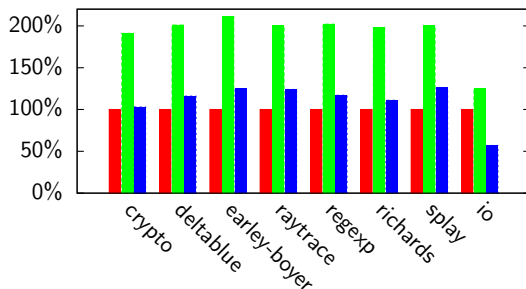
Normal Execution 
Serial Multi-Execution 
Parallel Multi-Execution 



Experimental Results: Execution Time

Normal Execution 
Serial Multi-Execution 
Parallel Multi-Execution 

Execution Time



Automatic parallelization

```
1 for (var i = 0; i < 100; ++i) {
2   var test = 0;
3   for (var j = 0; j < 10000; ++j) {
4     test += j;
5   }
6   if (i % 10 == 0) {
7     var hi_in = hi_input(); ← Latency
8     var lo_in = lo_input(); ← Latency
9     lo_output("#" + i + ". lo_in: '"
10      + lo_in + "'. hi_in is: '"
11      + hi_in + "'");
12     hi_output("#" + i + ". hi_in: '" ← Latency
13      + hi_in + "'. lo_in is: '"
14      + lo_in + "'");
15   }
16 }
```

Automatic parallelization

```

1 for (var i = 0; i < 100; ++i) {
2   var test = 0;
3   for (var j = 0; j < 10000; ++j) {
4     test += j;
5   }
6   if (i % 10 == 0) {
7     var hi_in = hi_input(); ← L Skip H Latency
8     var lo_in = lo_input(); ← Latency Reuse
9     lo_output("#" + i + ". lo_in: '"
10      + lo_in + "' . hi_in is: '"
11      + hi_in + "'");
12    hi_output("#" + i + ". hi_in: '" ← Skip Latency
13      + hi_in + "' . lo_in is: '"
14      + lo_in + "'");
15  }
16 }

```

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The technique's merits

Advantages:

Very strong Soundness guarantee

Very general No fundamental issues with parallelism, exceptions or other language features

Good precision No change for (termination-sensitively) non-interferent programs

Acceptable imprecision Interferent executions are modified in acceptable way (intuitive, no formalisation...)

Dynamic Run-time assignment of I/O channels to security levels

Downsides:

Performance Acceptable for some use cases?

Thank you for your attention.

Any questions?